System Feature Description: Importing Refutations into the GAPT Framework PxTP Workshop, Manchester

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Introduction

Context: GAPT Framework

GAPT = General Architecture for Proofs and Theorems provides (in different stages of development):

- Languages
 (Typed Lambda Calculus, First Order Logic, Higher Order Logic)
- Calculi (various Sequent Calculi, Resolution)
- Algorithms (Unification, Matching, ...)
- Interactive theorem prover (TAP)
- Proof Transformations
 (Proof Skolemization, Cut-elimination by Resolution, Herbrand Sequent Extraction)

CERES

Short overview of the CERES method

- Preprocessing of the input Sequent Calculus proof (Skolemization, Regularization)
- Extraction of the characteristic clause set
- Refutation of the characteristic clause set by an external resolution theorem prover
- Constructing proof projections to clauses from the characteristic clause set
- Constructing a proof in atomic-cut normal form from the refutation and the projections

Importing a Proof

Problems with Proof Parsing

- Variable renaming
- Substitutions not given
- Variety of inference rules
- Contraction of several inferences into one
- Incomplete or outdated documentation of the inference rules

Clause set:
$$\{ \vdash P(a); P(x) \vdash P(f(x)); \vdash f(x) = g(x); P(g(a)) \vdash \}$$

Refutation:

$$\cfrac{ \vdash P(a) \qquad P(x) \vdash P(f(x))}{ \vdash P(f(a))} \; \textit{Res} \; \sigma = \{x \mapsto a\} \qquad \qquad \vdash f(x) = g(x) \qquad \qquad \vdash p(g(a)) \vdash \qquad \qquad \qquad \vdash P(g(a)) \vdash \vdash$$

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Prover9:
```

```
1 f(x) = g(x). [assumption].
2 -P(x) | P(f(x)). [assumption].
3 -P(x) | P(g(x)). [copy(2),rewrite([1(2)])].
4 P(a). [assumption].
5 -P(g(a)). [assumption].
6 P(g(a)). [hyper(3,a,4,a)].
7 $F. [resolve(6,a,5,a)].
```

```
SPASS:
```

```
1[0:Inp] || -> equal(g(U),f(U))**.
2[0:Inp] P(U) || -> P(f(U))*.
3[0:Inp] || -> P(a)*.
4[0:Inp] || P(g(a))* -> .
5[0:Rew:1.0,4.0] || P(f(a))* -> .
7[0:Res:2.1,5.0] P(a) || -> .
8[0:SSi:7.0,3.0] || -> .
```

Vampire:

```
7. $false (1:0) [subsumption resolution 6,3]
```

- 3. 'P'(a) (0:2) [input]
- 6. "'P'(a) (1:2) [resolution 5,4]
- 4. ~'P'(g(a)) (0:3) [input]
- 5. 'P'(g(X0)) | ~'P'(X0) (0:5) [definition unfolding 2,1]
- 1. f(X0) = g(X0) (0:5) [input]
- 2. 'P'(f(X0)) | ~'P'(X0) (0:5) [input]

Vampire TPTP output:

```
fof(f7,plain,($false),
  inference(subsumption_resolution,[],[f6,f3])).
fof(f3,axiom,('P'(a)),
  file('simple.tptp',unknown)).
fof(f6,plain,(~'P'(a)),
  inference(resolution, [], [f5,f4])).
fof(f4,axiom,(~'P'(g(a))),
  file('simple.tptp',unknown)).
fof(f5,plain,((! [X0]: ('P'(g(X0)) | ~'P'(X0)))),
  inference(definition_unfolding,[],[f2,f1])).
fof(f1,axiom,((![X0]:(f(X0)=g(X0)))),
  file('simple.tptp',unknown)).
fof(f2,axiom,((![X0]:('P'(f(X0))|~'P'(X0)))).
  file('simple.tptp',unknown)).
```

Common Structure

• Inference label by: clause id, premise ids, clause, rule name

Problem

- Parse proof of an external resolution prover
- Fill in missing information
- Normalize proof to use only resolution and paramodulation

Approach

- Extract premises and target clause from proof step
- Use internal prover TAP to reprove each single step (forward resolution)
- Construct full refutation from the steps

The TAP Prover

- Simple resolution prover
- Intended for interactive use and experiments
- Commands based

TAP Internals

- Configuration: State + Commands Queue + Data
- State: clause set + guidance map
- Command: Function from configuration to list of successor states (possibly empty)
- Data: information passed only to following command, not stored in state

Commands for original use (interactive theorem prover)

- Resolve
- Paramodulation
- Factor
- Variants
- DeterministicAnd
- SetStream
- SetTargetClause
- InsertResolvent
- RefutationReached

Changes for Replaying

Changes

- Store resolution proofs instead of clauses
- Add new commands: Prover9Init, Replay, guidance commands

Prover9Init

- Pass clause set to theorem prover and parse result
- Schedule InsertResolvent and AddGuidedInitialClause command for each assumption
- Schedule Factor command for each factoring inference
- Schedule Replay command for every other inference step

Changes for Replaying

Replay

- Create new TAP instance
- Get proofs for premise clauses from guidance map
- Schedule SetClauseWithProof command for the premise clauses
- Schedule SetTargetClause command for the target clause
- Initialize prover to use Resolution and Paramodulation for proof search
- Start proof search
- Add proof found to guidance map and schedule InsertResolvent command for proof of target clause

Changes for Replaying

Guidance Map Management

- SetClauseWithProof
- AddGuidedInitialClause
- AddGuidedClauses
- GetGuidedClauses
- IsGuidedNotFound

Command Queue after Prover9Init

```
AddGuidedInitialClause(1, List(= (f(x), g(x))))
InsertResolvent
AddGuidedInitialClause(2, List(\neg P(x), P(f(x))))
InsertResolvent
```

Replay(List(0, 2, 1))

AddGuidedInitialClause(4, List(P(a)))

InsertResolvent

AddGuidedInitialClause(5, List($\neg P(g(a))$))

InsertResolvent

Replay(List(0, 3, 4))

Replay(List(0, 6, 5))



Replayed Example

$$\begin{array}{c} \frac{ \left| \begin{array}{c} \vdash f(x) = g(x) \\ \hline \vdash f(x_{\mathbf{6}}) = g(x_{\mathbf{6}}) \end{array} \right| Variant}{ \left| \begin{array}{c} \vdash P(x) \vdash P(f(x)) \\ \hline P(x_{\mathbf{5}}) \vdash P(f(x_{\mathbf{5}})) \end{array} \right| Variant} \\ \frac{ P(x_{\mathbf{5}}) \vdash P(g(x_{\mathbf{5}})) \\ \hline P(x_{\mathbf{10}}) \vdash P(g(x_{\mathbf{10}})) \end{array}}{ \left| \begin{array}{c} \vdash P(g(a)) \vdash \\ \hline P(g(a)) \vdash \end{array} \right|} \begin{array}{c} Variant \\ Variant \\ Res \ \sigma = \{x_{\mathbf{10}} \mapsto a\} \end{array} \\ \begin{array}{c} \vdash P(g(a)) \vdash \\ Res \end{array}$$

Pitfalls

Pitfalls

- Forward reasoning prevents some strategies
 - Factorization can not only be applied after an inference step
 - No reflexivity rule: add reflexivity axiom or unfold rule
 - Equations might get flipped
 - Expectation that a single inference is provable in few steps not met

Conclusion and Future Work

- Normalized proof with instantiations needed for cut-elimination and Herbrand sequent extraction
- Replay of Prover9 proofs works for small examples, performance issues for larger ones
- Macro rules with large numbers of premises need specialized handling (necessary for Vampire/SPASS/E/... integration)

Thanks for the attention!